

# Image Segmentation by Image Foresting Transform with Non-smooth Connectivity Functions

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**Abstract**—In the framework of the Image Foresting Transform (IFT), there is a class of connectivity functions that were vaguely explored, which corresponds to the non-smooth connectivity functions (NSCF). These functions are more adaptive to cope with the problems of field inhomogeneity, which are common in MR images of 3 Tesla. In this work, we investigate the NSCF from the standpoint of theoretical and experimental aspects. We formally classify several non-smooth functions according to a proposed diagram representation. Then, we investigate some theoretical properties for some specific regions of the diagram. Our analysis reveals that many NSCFs are, in fact, the result of a sequence of optimizations, each of them involving a maximal set of elements, in a well-structured way. Our experimental results indicate that substantial improvements can be obtained by NSCFs in the 3D segmentation of MR images of 3 Tesla, when compared to smooth connectivity functions.

**Keywords**—graph search algorithms; image foresting transform; non-smooth connectivity functions

## I. INTRODUCTION

Quantitative analysis of brain structures from magnetic resonance images (MRI) has played an important role in research in neurology and can be very useful in the diagnosis and treatment of diseases related to changes in the anatomy of the human brain [1]. Analysis techniques, however, require a precise definition of the extent of the three-dimensional structures under study. This operation called image segmentation is one of the most fundamental and challenging problems in image processing and computer vision [2]. In medical image analysis, accurate segmentation results commonly require the user intervention because of the presence of structures with ill-defined borders, intensity non-standardness among images, field inhomogeneity, noise, artifacts, partial volume effects, and their interplay [1], [3].

Methods based on the Image Foresting Transform (IFT) [4] have been successfully used in the segmentation of 1.5 Tesla MR datasets [5], [6], [7]. However, inhomogeneity effects are stronger in higher magnetic fields, and it is extremely important to define the optimal solution for these images. In the framework of the IFT, there is a class of connectivity functions that were little explored, which corresponds to the non-smooth connectivity functions (NSCF). These functions are more adaptive to cope with the problems of field inhomogeneity, which are common in MR images of 3 Tesla [8].

Their solutions could also be useful in the handling of natural images with strong variation in brightness/shadow.

Historically, the usage of non-smooth connectivity functions (NSCF) has been avoided, because their optimality, in terms of the paths computed by the IFT, may not be guaranteed according to [4]. Early studies only considered the usage of NSCFs, such as  $f_{euc}$  (Eq. 3 with 8-connected adjacency), to compute an approximation of the Euclidean distance (EDT) [9], [10], although some other studies by Herman et al. have suggested some practical advantages of using NSCF in the context of image segmentation [11], [12]. More recently, Strand et al. proposed the *Minimum Barrier Distance* (MBD) [13], which, in digital setting, leads to a non-smooth connectivity function. In empirical comparisons, a favorable outcome has been reported for MBD [13]<sup>1</sup>. Also, in other recent works, it has been proven that some non-smooth connectivity functions can lead to optimum results according to other optimality criteria, such as a graph-cut measure [14], [15], or an energy function for optimum boundary tracking [16]. This brought new lights to the research involving NSCF.

In this work, we investigate the NSCF from the standpoint of theoretical and experimental aspects. We formally classify several non-smooth functions according to a proposed diagram representation. Then, we investigate some theoretical properties for some specific regions of the diagram. Our analysis reveals that many NSCFs are, in fact, the result of a sequence of optimizations, each of them involving a maximal set of elements, in a well-structured way. This process resembles the original definition of *Iterative Relative Fuzzy Connectedness* (IRFC) [17], which was defined as the result of a sequence of optimizations by *Relative Fuzzy Connectedness* (RFC) [18]. Our experimental results, using a robot user [19], indicate that substantial improvements can be obtained by NSCFs in the segmentation of MR images of 3 Tesla, when compared to smooth connectivity functions, and their related methods [17], [20].

For the sake of completeness in presentation, Section II includes the relevant previous work of image segmentation by IFT. Our analysis of non-smooth connectivity functions

<sup>1</sup>In [13], an algorithm that finds an approximation of the *Minimum Barrier Distance* was considered, by computing two IFTs separately, one for each term in Eq. 5.

is presented in Section III. In Section IV, we evaluate the NSCFs in the segmentation of MR images of 3 Tesla, and our conclusions are stated in Section V.

## II. IMAGE FORESTING TRANSFORM

An image can be interpreted as a graph  $G = (\mathcal{I}, \mathcal{A})$  whose nodes are the image pixels in its image domain  $\mathcal{I} \subset Z^N$ , and whose arcs are the pixel pairs  $(s, t)$  in  $\mathcal{A}$  (e.g., 4-neighborhood, or 8-neighborhood, in case of 2D images, and 6-neighbors in 3D). The *adjacency relation*  $\mathcal{A}$  is a binary relation on  $\mathcal{I}$ . We use  $t \in \mathcal{A}(s)$  and  $(s, t) \in \mathcal{A}$  to indicate that  $t$  is adjacent to  $s$ . Each arc  $(s, t) \in \mathcal{A}$  may have a weight  $w(s, t) \geq 0$ , such as a dissimilarity measure between pixels  $s$  and  $t$  (e.g.,  $w(s, t) = |I(t) - I(s)|$  for a single channel image with values given by  $I(t)$ ).

For a given image graph  $G = (\mathcal{I}, \mathcal{A})$ , a path  $\pi_t = \langle t_1, t_2, \dots, t_n = t \rangle$  is a sequence of adjacent pixels with terminus at a pixel  $t$ . A path is *trivial* when  $\pi_t = \langle t \rangle$ . A path  $\pi_t = \pi_s \cdot \langle s, t \rangle$  indicates the extension of a path  $\pi_s$  by an arc  $(s, t)$ . A *predecessor map* is a function  $P$  that assigns to each pixel  $t$  in  $\mathcal{I}$  either some other adjacent pixel in  $\mathcal{I}$ , or a distinctive marker *nil* not in  $\mathcal{I}$  — in which case  $t$  is said to be a *root* of the map. A *spanning forest* is a predecessor map which contains no cycles — i.e., one which takes every pixel to *nil* in a finite number of iterations. For any pixel  $t \in \mathcal{I}$ , a spanning forest  $P$  defines an induced path  $\pi_t^P$  recursively as  $\langle t \rangle$  if  $P(t) = \text{nil}$ , and  $\pi_s^P \cdot \langle s, t \rangle$  if  $P(t) = s \neq \text{nil}$ .

A *connectivity function* computes a value  $f(\pi_t)$  for any path  $\pi_t$ , usually based on arc weights. A path  $\pi_t$  is *optimum* if  $f(\pi_t) \leq f(\tau_t)$  for any other path  $\tau_t$  in  $G$ . By taking to each pixel  $t \in \mathcal{I}$  one optimum path with terminus  $t$ , we obtain the optimum-path value  $V(t)$ , which is uniquely defined by  $V(t) = \min_{\forall \pi_t \text{ in } G} \{f(\pi_t)\}$ . The *image foresting transform* (IFT) [4] takes an image graph  $G = (\mathcal{I}, \mathcal{A})$ , and a path-value function  $f$ ; and assigns one optimum path  $\pi_t$  to every pixel  $t \in \mathcal{I}$  such that an *optimum-path forest*  $P$  is obtained — i.e., a spanning forest where all induced paths are optimum. However,  $f$  must be *smooth* (Definition 1), otherwise, the paths may not be optimum, as demonstrated in [4].

**Definition 1** (Smooth path-value function). A *path-value function*  $f$  is smooth if for any node  $t \in \mathcal{I}$ , there is a simple optimum path  $\pi_t$  (no repeated vertices), which either is trivial, or has the form  $\pi_s \cdot \langle s, t \rangle$  where

$$(C1) \quad f(\pi_s) \leq f(\pi_t),$$

$$(C2) \quad \pi_s \text{ is optimum,}$$

$$(C3) \quad \pi_s \text{ is optimum, and for any optimum path } \pi'_s, f(\pi'_s \cdot \langle s, t \rangle) = f(\pi_t)^2.$$

Any path  $\pi_t^P$  encoded in an optimum-path forest  $P$  (i.e., the paths starting in a root node) is a complete-optimum path, according to the following definition:

**Definition 2** (Complete-optimum path). A path  $\pi_{t_n} = \langle t_1, t_2, \dots, t_n \rangle$  is complete optimum if all paths  $\pi_{t_i} =$

$\langle t_1, t_2, \dots, t_i \rangle$ ,  $i = 1, 2, \dots, n$  are optimum paths.

The cost of a trivial path  $\pi_t = \langle t \rangle$  is usually given by a handicap value  $H(t)$ , while the connectivity functions for non-trivial paths follow a path-extension rule. For example:

$$f_{\max}(\pi_s \cdot \langle s, t \rangle) = \max\{f_{\max}(\pi_s), w(s, t)\} \quad (1)$$

$$f_{\Sigma}(\pi_s \cdot \langle s, t \rangle) = f_{\Sigma}(\pi_s) + w(s, t) \quad (2)$$

$$f_{\text{euc}}(\pi_s \cdot \langle s, t \rangle) = \|t - R(\pi_s)\|^2 \quad (3)$$

where  $w(s, t) \geq 0$  is a fixed arc weight, and  $R(\pi_t)$  is the origin/root of a path  $\pi_t$ . The function  $f_{\text{euc}}$  may not be smooth, depending on the adjacency  $\mathcal{A}$  [4].

We consider image segmentation from two seed sets,  $\mathcal{S}_o$  and  $\mathcal{S}_b$  ( $\mathcal{S}_o \cap \mathcal{S}_b = \emptyset$ ), containing pixels selected inside and outside the object, respectively. The search for optimum paths is constrained to start in  $\mathcal{S} = \mathcal{S}_o \cup \mathcal{S}_b$  (i.e.,  $H(t) = 0$  for all  $t \in \mathcal{S}$ , and  $H(t) = +\infty$  otherwise). The image is partitioned into two optimum-path forests — one rooted at the internal seeds, defining the object, and the other rooted at the external seeds, representing the background [21]. A label,  $L(t) = 1$  for all  $t \in \mathcal{S}_o$  and  $L(t) = 0$  for all  $t \in \mathcal{S}_b$ , is propagated to all unlabeled pixels during the computation [4].

### A. General IFT Algorithm

Algorithm 1 obtains an optimum-path forest  $P$  by minimizing a smooth path-value function  $f$ , or a spanning forest  $P$  in the case of a non-smooth function.

#### Algorithm 1. – GENERAL IFT ALGORITHM

INPUT: Image graph  $(\mathcal{I}, \mathcal{A})$ , seed sets  $\mathcal{S}_o$  and  $\mathcal{S}_b$ , and path-value function  $f$ .  
 OUTPUT: Optimum-path forest  $P$ , the minimum path-value map  $V$  and label map  $L$ .  
 AUXILIARY: Priority queue  $Q$ , variable *tmp*, and an array of *status*.

1. **For each**  $t \in \mathcal{S}_o$ , **do**  $L(t) \leftarrow 1$ .
2. **For each**  $t \in \mathcal{S}_b$ , **do**  $L(t) \leftarrow 0$ .
3. **For each**  $t \in \mathcal{I}$ , **do**
4.     Set  $P(t) \leftarrow \text{nil}$  and  $V(t) \leftarrow f(\langle t \rangle)$ .
5.     Set  $\text{status}(t) \leftarrow 0$ .
6.     **If**  $V(t) \neq +\infty$ , **then insert**  $t$  in  $Q$ .
7. **While**  $Q \neq \emptyset$ , **do**
8.     Remove  $s$  from  $Q$  such that  $V(s)$  is minimum.
9.     Set  $\text{status}(s) \leftarrow 1$ .
10.     **For each**  $t \in \mathcal{A}(s)$ , such that  $\text{status}(t) = 0$ , **do**
11.         Compute  $\text{tmp} \leftarrow f(\pi_s^P \cdot \langle s, t \rangle)$ .
12.         **If**  $\text{tmp} < V(t)$ , **then**
13.             **If**  $V(t) \neq +\infty$ , **then remove**  $t$  from  $Q$ .
14.             Set  $P(t) \leftarrow s$ ,  $V(t) \leftarrow \text{tmp}$ .
15.             Set  $L(t) \leftarrow L(s)$  and insert  $t$  in  $Q$ .

The “for” loops initialize maps and insert pixels with finite trivial-path values in  $Q$  (Lines 1–6). The minima of the initial map  $V$  compete with each other and some of them become roots of the forest. The main loop computes optimum paths from the minima to every pixel  $s$  in a non-decreasing order of value (Lines 7–15). At each iteration, a path  $\pi_s^P$  of minimum value  $V(s)$  is obtained in  $P$  when we remove its last pixel  $s$  from  $Q$  (Line 8). The rest of the lines evaluate if the path

<sup>2</sup>Our condition C3 corresponds, in fact, to the intersection of the original conditions C2 and C3 from [4], in order to keep the diagram easier.

$\pi_s^P \cdot \langle s, t \rangle$  that reaches an adjacent pixel  $t$  through  $s$  is cheaper than the current path  $\pi_t^P$  in  $P$  and update  $Q$ ,  $V(t)$ ,  $L(t)$ , and  $P(t)$  accordingly.

### III. NON-SMOOTH CONNECTIVITY FUNCTIONS

Clearly, from Definition 1, we have that a connectivity function is not smooth if it doesn't satisfy at least one of the conditions C1, C2 or C3. Note that, if a function  $f$  does not satisfy C2, this implies that the condition C3 is not satisfied either. In order to further distinguish between different classes of non-smooth functions, we also evaluate the functions with respect to the condition C4, which is defined below:

**Definition 3** (Condition C4). *A path-value function  $f$  satisfies the condition C4, if for any node  $s \in \mathcal{I}$  the following condition is verified  $\forall t \in \mathcal{A}(s)$ :*

- For any paths  $\pi_s$  and  $\pi'_s$  ending at  $s$ , if  $f(\pi_s) = f(\pi'_s)$ , then we have  $f(\pi_s \cdot \langle s, t \rangle) = f(\pi'_s \cdot \langle s, t \rangle)$ .

For a general image graph, the classification of various non-smooth functions into the sets C1, C2, C3, and C4 (such that a function  $f$  is in a set  $C_i$  if and only if it satisfies the condition  $C_i$ ) is shown in the proposed diagram illustrated in Figure 1. The functions  $f_{\max}^{bkg}$ ,  $f_\omega$ ,  $f_\omega^\circ$ ,  $f_{i,\omega}$ , and  $f_{o,\omega}$  were studied in [16], [14], [15]<sup>3</sup>. The other functions (i.e.,  $f_{\Sigma \max}$ ,  $f_{\uparrow}$ ,  $f_{\max|\Delta I|}$ ,  $f_{\Sigma|\Delta I|}$ ,  $f_I$ , and  $f_{\Sigma \max}^{lex}$ ) are defined in the Appendix section, including a proof of their respective positions in the proposed diagram.

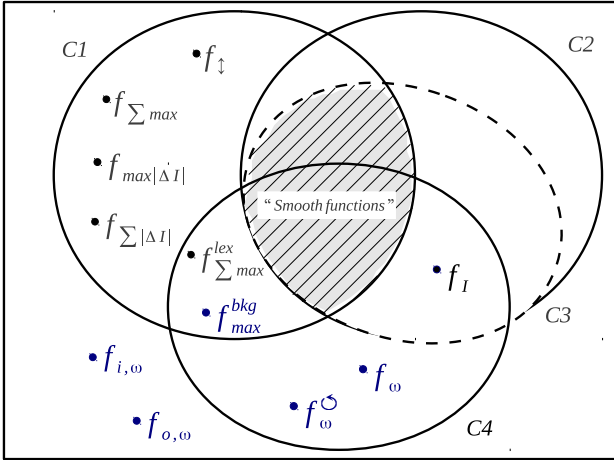


Fig. 1. Schematic representation of the relations between smooth and non-smooth connectivity functions: C1, C2, C3, and C4 are sets of connectivity functions that satisfy these respective conditions for a general graph.

The next propositions reveal some interesting theoretical properties of any non-smooth connectivity function  $f \in (C1 \cap C4) \setminus C2$ .

**Proposition 1.** *Consider a connectivity function  $f$ , classified as  $f \in (C1 \cap C4) \setminus C2$  for a general graph. If for a given image*

<sup>3</sup>The proof of their positions in the diagram were omitted due to the lack of space.

graph  $G = (\mathcal{I}, \mathcal{A})$  and set of seeds  $\mathcal{S}$ , there exists an optimum-path forest for  $f$ , then any spanning forest  $P$  computed in  $G$  by Algorithm 1 for  $f$  is an optimum-path forest.

*Proof:* Although, in the general case, the connectivity function  $f$  violates C2, for this particular graph,  $f$  actually satisfies the condition C2. Note that, for all  $t \in \mathcal{I}$  we can take an optimum path  $\pi_t = \pi_s \cdot \langle s, t \rangle$  where  $\pi_s$  is optimum from the existing optimum-path forest. Hence,  $f$  behaves as a smooth function for this particular graph, leading to optimum results. ■

When, on previous works, it is said that Algorithm 1 computes a spanning forest that may not be optimum for non-smooth functions, you might think that there was such an optimum-path forest, but the algorithm was not able to find it. However, as a consequence of Proposition 1, we have that for any connectivity function  $f \in (C1 \cap C4) \setminus C2$ , Algorithm 1 only fails to find an optimum-path forest, if no such forest exists. Thus, it is not a fault of the algorithm itself, but a fault of the problem specification which has no valid solution in the form of an optimum-path forest.

**Proposition 2.** *Consider a function  $f \in (C1 \cap C4) \setminus C2$ . For a given image graph  $G = (\mathcal{I}, \mathcal{A})$ , and set of seeds  $\mathcal{S}$ , let  $\mathcal{O}$  be the set of all pixels  $t \in \mathcal{I}$ , such that there exists a complete-optimum path  $\pi_t$  for  $f$ . In any spanning forest  $P$  computed in  $G$  by Algorithm 1 for  $f$ , all the paths  $\tau_t^P$  with  $t \in \mathcal{O}$  are optimum paths.*

*Proof:* Let  $\pi_t = \langle t_1, t_2, \dots, t_i, \dots, t_n = t \rangle$  be a given complete-optimum path for  $f$ , and  $\tau_t^P = \langle s_1, s_2, \dots, s_m = t \rangle$  be the computed path in the spanning forest  $P$ . We have the following proof by mathematical induction:

- **The basis:** Show that the statement holds for  $n = 1$ . In this case, we have  $t = t_1$ , and therefore  $\pi_t = \langle t_1 \rangle$ . Since  $f(\langle t_1 \rangle)$  is a finite value,  $t_1$  is inserted into the priority queue  $Q$  (Line 6 of Alg. 1). Since  $f(\langle t_1 \rangle)$  is optimum, we can safely assure that it won't leave  $Q$  (Line 8) with a cost  $V(t_1)$  worse than  $f(\langle t_1 \rangle)$ , therefore, the computed path  $\tau_t^P$  must be optimum.
- **The inductive step:** Assume that the statement is true for the complete-optimum path  $\pi_{t_i}$ ,  $i \geq 1$ . We must prove that it also holds for the complete-optimum path  $\pi_{t_{i+1}} = \pi_{t_i} \cdot \langle t_i, t_{i+1} \rangle$ . By the hypothesis we have that the path  $\tau_{t_i}^P$  in  $P$  is optimum. Given that the condition C4 is satisfied ( $f \in (C1 \cap C4) \setminus C2$ ), we may conclude that  $f(\tau_{t_i}^P \cdot \langle t_i, t_{i+1} \rangle) = f(\pi_{t_i} \cdot \langle t_i, t_{i+1} \rangle) = f(\pi_{t_{i+1}})$ , so we may conclude that  $\tau_{t_i}^P \cdot \langle t_i, t_{i+1} \rangle$  is optimum. By algorithm construction, if  $t_{i+1}$  was not yet conquered, we have that this extended path  $\tau_{t_i}^P \cdot \langle t_i, t_{i+1} \rangle$  will be evaluated by the Algorithm 1 (Lines 10-15), and since it is optimum, it won't be posteriorly replaced by any other path (as guaranteed by the strict inequality in Line 12). If  $t_{i+1}$  was already conquered by following a different path in  $P$ , its cost cannot be worse than  $f(\pi_{t_{i+1}})$ , because  $f \in C1$ . So in the end,  $\tau_{t_{i+1}}^P$  will be an optimum path.

The successive application of the previous proposition gives a characterization of Algorithm 1 for  $f \in (C1 \cap C4) \setminus C2$ , as the result of a sequence of optimizations, where each optimization step involves a maximal set of elements, in a well-structured way.

Consider the following definitions: Let  $\mathcal{E}^{set}(X) = \{\forall(s,t) \in \mathcal{A} \mid s \in X \wedge t \in X\}$  denote the set of all arcs interconnecting nodes in the set  $X$ ,  $\mathcal{E}^{path}(\pi)$  denote the set of all arcs in the path  $\pi$  (i.e.,  $\mathcal{E}^{path}(\pi) = \{\forall(t_i, t_{i+1}) \text{ for } 1 \leq i \leq k-1 \mid \pi = \langle t_1, t_2, \dots, t_i, t_{i+1}, \dots, t_k \rangle\}$ ),  $\mathcal{E}^{cut}(X, Y) = \{\forall(s,t) \in \mathcal{A} \mid s \in X \wedge t \in Y\}$ ,  $\mathcal{E}^{pred}(X) = \bigcup_{t \in X} \mathcal{E}^{path}(\pi_t^P)$ .

In the first optimization step, optimum paths  $\tau_t^P$  are computed for all  $t \in \mathcal{O}$  (Proposition 2). Let's denote  $\mathcal{O}$  as  $\mathcal{O}^1$  for this first step. In the next optimization step, consider the sub-graph  $G^2 = (\mathcal{I}, \mathcal{E}^{set}(\mathcal{I} \setminus \mathcal{O}^1) \cup \mathcal{E}^{cut}(\mathcal{O}^1, \mathcal{I} \setminus \mathcal{O}^1) \cup \mathcal{E}^{pred}(\mathcal{O}^1))$ . A second path optimization is performed, by computing a second IFT, but now in  $G^2$ <sup>4</sup>. Since the arcs interconnecting nodes in  $\mathcal{O}^1$ , are reduced to the arcs in the previous forest  $P$  (i.e.,  $\mathcal{E}^{pred}(\mathcal{O}^1)$ ) in  $G^2$ , we have that the optimum paths  $\tau_t^P$ , computed on the previous step, will remain optimum in the new graph  $G^2$ . So the optimum paths  $\tau_t^P$  with  $t \in \mathcal{O}^1$  will start a new competition, seeking for their best extensions to the other pixels in  $\mathcal{I} \setminus \mathcal{O}^1$ . By applying the Proposition 2 on this new optimization problem one more time, we have that this second IFT will conquer a new maximal set of pixels  $\mathcal{O}^1 \cup \mathcal{O}^2$  that can be reached by optimum paths in  $G^2$ . We can then repeat this process over again. The condition  $C1^5$  guarantees that at least one new element will be conquered at each step, so that this process will repeat until  $\bigcup_{v_i} \mathcal{O}^i = \mathcal{I}$ .

#### IV. EXPERIMENTS

In this section, we present the quantitative analysis of the accuracy for segmenting the brain in 3D, in experiments involving a 3 Tesla MR-image dataset, containing ten T1-weighted 3D images. The image dataset included the head and, at least, a small portion of the neck of male and female adults with normal brains. A specialist delineated the ground-truths (GT)s. The seeds were chosen by a robot user, the method introduced by Gulshan et al. [19], to simulate user interaction of interactive segmentation by placing brush strokes automatically to iteratively, and interactively, complete the segmentation task. The procedure, including the choice of seeds (based on the knowledge of the ground truth) and the successive delineations, is iterative. The initial seeds for the object are placed at the point(s) in the object farthest from the boundary; similarly, for the background. At each successive iteration, the robot always places a spherical brush stroke in the largest connected component of the segmentation error area (placed at the point(s) farthest from the boundary of the component), and updates the segmentation. The process is

<sup>4</sup> By Algorithm 1, the pixels  $s \in \mathcal{O}^1$  have  $status(s) = 1$ , and their paths can't be changed. Therefore, we consider a new graph  $G^2$ , where the edges interconnecting nodes in  $\mathcal{O}^1$  which are not in the forest  $P$  are disregarded.

<sup>5</sup>By the hypothesis, we know that  $f$  satisfies the condition  $C1$ .

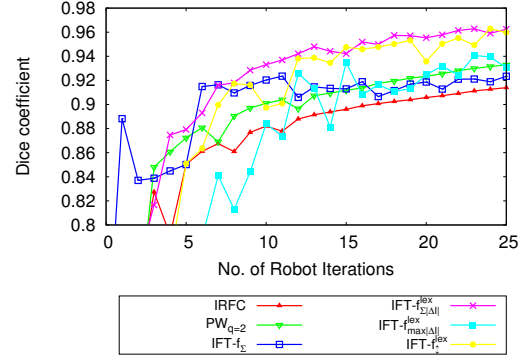


Fig. 2. Results using a robot user for segmenting the brain dataset.

repeated up to 25 times, generating a sequence of 25 simulated user strokes.

We used a spherical brush with a diameter of 10 voxels. Figure 2 shows the experimental curves, where *IRFC* and  $PW_{q=2}$  represent different algorithms related to the smooth function  $f_{\max}$  [17], [20], and we used  $w(s,t) = G(s) + G(t)$ , where  $G(s)$  is the magnitude of Sobel gradient at a voxel  $s$ . Clearly,  $f_{\Sigma|\Delta I|}^{lex}$  presented the best accuracy. Figure 3 shows one example for user-selected markers. These results emphasize the importance of non-smooth connectivity functions. The non-smooth connectivity function  $f_{\Sigma|\Delta I|}^{lex}$  is a variation of  $f_{\Sigma|\Delta I|}$  (Equation 9), in order to guarantee that  $f_{\Sigma|\Delta I|}^{lex} \in (C1 \cap C4) \setminus C2$ . This transformation is similar to the one used to get  $f_{\Sigma_{\max}}^{lex} \in (C1 \cap C4) \setminus C2$  from  $f_{\Sigma_{\max}}$  (see Appendix section). The function  $f_{\Sigma|\Delta I|}^{lex}$  gives pairs of values that should be compared according to the lexicographical order. The first component is the non-smooth function  $f_{\Sigma|\Delta I|}$  (Equation 9), and the second is the priority level of the seed/root for that path. The lower its value the higher is its priority. In interactive segmentation, we give lower priority for new inserted seeds, since they are used mainly for corrective actions, so that we can keep their effects more locally. The same process was done for  $f_{\max|\Delta I|}^{lex} \in (C1 \cap C4) \setminus C2$  and  $f_{\downarrow}^{lex} \in (C1 \cap C4) \setminus C2$ , in relation to  $f_{\max|\Delta I|}$  (Eq. 8) and  $f_{\downarrow}$  (Eq. 5), respectively.

Although the results obtained by the smooth function  $f_{\max}$  (e.g., *IRFC*) could be improved by using a better gradient-like image, this would require an undesirable training step (implicitly or explicitly) in order to enhance the object boundary [22]. IFT using  $f_{\Sigma|\Delta I|}^{lex}$  not only avoids any training stage, but it can also better handle the inhomogeneity effects (Figure 4).

#### V. CONCLUSION

We have shown a promising subset of non-smooth connectivity functions (i.e.,  $f \in (C1 \cap C4) \setminus C2$ ), including a theoretical analysis of its properties. In particular, the function  $f_{\Sigma|\Delta I|}^{lex}$  demonstrated very good results for the 3D segmentation of MR images of 3 Tesla. We intend to combine it with statistical models for automatic segmentation, and to extend this analysis to other non-smooth connectivity functions.

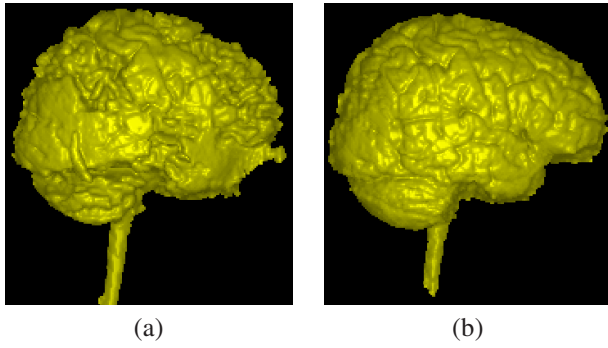


Fig. 3. Brain segmentation results for the same user-selected markers by (a)  $f_{\max}$ , and (b)  $f_{\sum |\Delta I|}^{\text{lex}}$ .

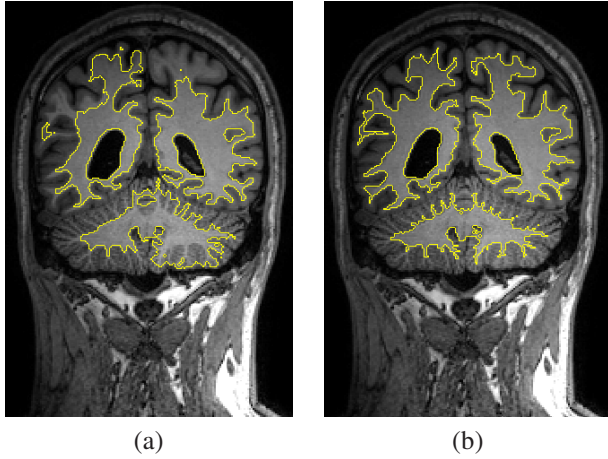


Fig. 4. White matter segmentation in 3D for the same user-selected markers. Results by (a)  $f_{\max}$  over an enhanced gradient, and (b)  $f_{\sum |\Delta I|}^{\text{lex}}$ .

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#### APPENDIX A

##### A. Function $f_{\sum \max}$ (sum of $f_{\max}$ )

Let  $f_{\sum \max}$  be the function defined by:

$$f_{\sum \max}(\pi_t = \langle t \rangle) = \begin{cases} 0, & \text{if } t \in S, \\ +\infty, & \text{otherwise.} \end{cases}$$

$$f_{\sum \max}(\pi_t = \pi_s \cdot \langle s, t \rangle) = f_{\sum \max}(\pi_s) + f_{\max}(\pi_t), \quad (4)$$

where  $S$  is a set of seeds.

From the diagram, we have that  $f_{\sum \max} \in C^1 \setminus (C^2 \cup C^4)$ .

##### Proof:

- C1. There is an optimum path  $\pi_t = \pi_s \cdot \langle s, t \rangle$ , such that  $f_{\sum \max}(\pi_s) \leq f_{\sum \max}(\pi_t)$ . For  $\omega(s, t) \geq 0$  we have that  $f_{\max}(\pi_t) = \max\{f_{\max}(\pi_s), \omega(s, t)\} \geq 0$ , then  $f_{\sum \max}(\pi_s) \leq f_{\sum \max}(\pi_s) + f_{\max}(\pi_t)$ . So, from (4) we obtain  $f_{\sum \max}(\pi_s) \leq f_{\sum \max}(\pi_t)$ . Therefore,  $f_{\sum \max}$  satisfies this condition.
- C2. There is an optimum path  $\pi_t = \pi_s \cdot \langle s, t \rangle$ , where  $\pi_s$  is optimum.

The proof that  $f_{\Sigma \max}$  does not lie in C2 can be demonstrated by the following counterexample. In Fig. 5, for the optimum path  $\pi_t = \pi_s \cdot \langle s, t \rangle$ , we have that  $\pi_s$  is not an optimum path to  $s$ . There is another path  $\pi'_s$  that offers a better cost to  $s$  than  $\pi_s$  (i.e.  $f_{\Sigma \max}(\pi'_s) = 9 \leq f_{\Sigma \max}(\pi_s) = 10$ ). So the condition C2 may not be verified for  $f_{\Sigma \max}$ .

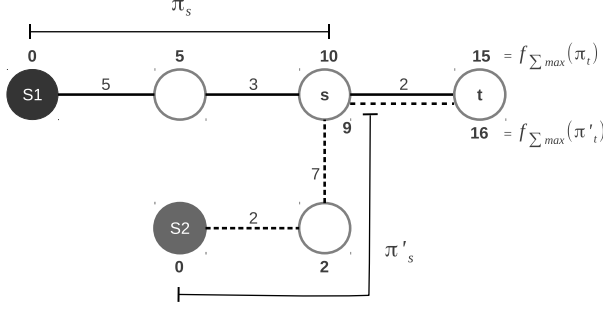


Fig. 5. Counterexample of the second and third conditions (C2 and C3) for the function  $f_{\Sigma \max}$ , where  $\pi_t = \pi_s \cdot \langle s, t \rangle$  is an optimum path, and  $\pi'_t = \pi'_s \cdot \langle s, t \rangle$  is a non-optimal path, from seeds  $s_1$  and  $s_2$ , respectively.

- C3. For any optimum path  $\pi'_s$  ending at  $s$ ,  $f_{\Sigma \max}(\pi'_s \cdot \langle s, t \rangle) = f_{\Sigma \max}(\pi_t)$ .  
 In Fig. 5 we can observe that  $f_{\Sigma \max}(\pi'_s \cdot \langle s, t \rangle) = 16 \neq f_{\Sigma \max}(\pi_t) = 15$ . Hence, C3 is violated by  $f_{\Sigma \max}$ .
- C4. For any paths  $\pi_s$  and  $\pi'_s$  ending at  $s$ , if  $f_{\Sigma \max}(\pi_s) = f_{\Sigma \max}(\pi'_s)$ , then we have  $f_{\Sigma \max}(\pi_s \cdot \langle s, t \rangle) = f_{\Sigma \max}(\pi'_s \cdot \langle s, t \rangle)$ .  
 In Fig. 6, for the optimum paths  $\pi_s$  and  $\pi'_s$  we have  $f_{\Sigma \max}(\pi_s \cdot \langle s, t \rangle) = 15 \neq f_{\Sigma \max}(\pi'_s \cdot \langle s, t \rangle) = 16$ . Therefore, clearly we have that  $f_{\Sigma \max}$  doesn't satisfy property C4.

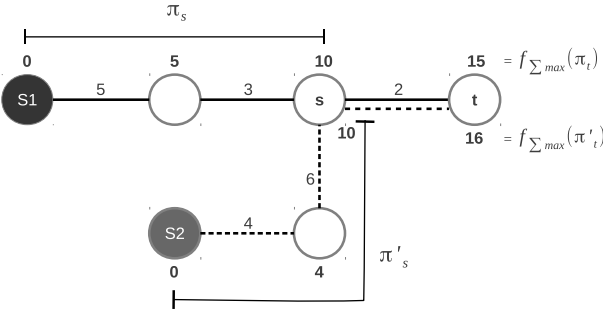


Fig. 6. Counterexample of the property C4 for the function  $f_{\Sigma \max}$ , where  $\pi_t = \pi_s \cdot \langle s, t \rangle$ ,  $\pi_s$ , and  $\pi'_s$  are optimum paths,  $\pi'_t = \pi'_s \cdot \langle s, t \rangle$  is a non-optimal path, and  $s_1, s_2$  are seeds.

### B. The barrier distance function

Let  $f_{\uparrow}$  be the function defined by:

$$f_{\uparrow}(\pi_t = \langle t \rangle) = \begin{cases} 0, & \text{if } t \in S, \\ +\infty, & \text{otherwise.} \end{cases}$$

$$f_{\uparrow}(\pi_t = \pi_s \cdot \langle s, t \rangle) = f_{I \max}(\pi_t) - f_{I \min}(\pi_t), \quad (5)$$

where  $S$  is a set of seeds, and  $f_{I \max}$  and  $f_{I \min}$  are functions that take the maximum and minimum intensity values along the path, respectively:

$$f_{I \max}(\pi_t = \langle t \rangle) = I(t)$$

$$f_{I \max}(\pi_t = \pi_s \cdot \langle s, t \rangle) = \max\{f_{I \max}(\pi_s), I(t)\}, \quad (6)$$

and

$$f_{I \min}(\pi_t = \langle t \rangle) = I(t)$$

$$f_{I \min}(\pi_t = \pi_s \cdot \langle s, t \rangle) = \min\{f_{I \min}(\pi_s), I(t)\}, \quad (7)$$

where  $I(t)$  is the intensity of a pixel  $t$ .

From the diagram, we have that  $f_{\uparrow} \in C1 \setminus (C2 \cup C4)$ .

**Proof:**

- C1.  $f_{\uparrow}(\pi_s) \leq f_{\uparrow}(\pi_t)$ .  
 Since  $f_{I \max}(\pi_s) \leq f_{I \max}(\pi_t)$ , and  $f_{I \min}(\pi_t) \leq f_{I \min}(\pi_s)$ , then,  $f_{I \max}(\pi_s) + f_{I \min}(\pi_t) \leq f_{I \max}(\pi_t) + f_{I \min}(\pi_s)$  and,  $f_{I \max}(\pi_s) - f_{I \min}(\pi_s) \leq f_{I \max}(\pi_t) - f_{I \min}(\pi_t)$ . So, from (5) we obtain  $f_{\uparrow}(\pi_s) \leq f_{\uparrow}(\pi_t)$ . Hence,  $f_{\uparrow}$  satisfies the condition C1.
- C2.  $\pi_s$  is optimum.  
 In Fig. 7, for the optimum path  $\pi_t = \pi_s \cdot \langle s, t \rangle$ , we have that  $\pi_s$  is not optimum, since there is another path  $\pi'_s$  which offers a better cost than  $\pi_s$  to  $s$  (i.e.,  $f_{\uparrow}(\pi'_s) = 6 \leq f_{\uparrow}(\pi_s) = 7$ ). Hence, C2 does not hold for  $f_{\uparrow}$ .

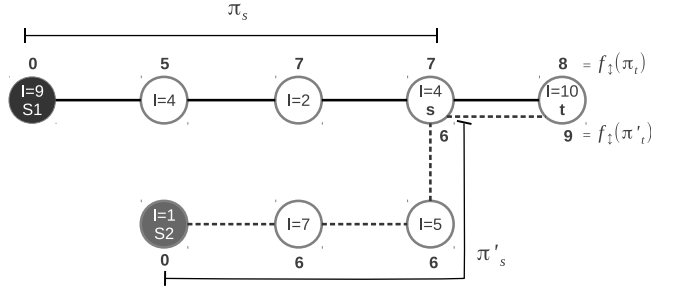


Fig. 7. Counterexample for the second and third conditions (C2 and C3) for the function  $f_{\uparrow}$ , where  $\pi_t = \pi_s \cdot \langle s, t \rangle$  is an optimum path, and  $\pi'_t = \pi'_s \cdot \langle s, t \rangle$  is a non-optimal path, from seeds  $s_1$  and  $s_2$ , respectively.

- C3. For any optimum path  $\pi'_s$ ,  $f_{\uparrow}(\pi'_s \cdot \langle s, t \rangle) = f_{\uparrow}(\pi_t)$ .  
 In Fig. 7 we have  $f_{\uparrow}(\pi'_s \cdot \langle s, t \rangle) = 9 \neq f_{\uparrow}(\pi_t = \pi_s \cdot \langle s, t \rangle) = 8$ . Hence, we have that  $f_{\uparrow} \notin C3$ .
- C4. For any paths  $\pi_s$  and  $\pi'_s$  ending at  $s$ , if  $f_{\uparrow}(\pi_s) = f_{\uparrow}(\pi'_s)$ , then  $f_{\uparrow}(\pi_s \cdot \langle s, t \rangle) = f_{\uparrow}(\pi'_s \cdot \langle s, t \rangle)$ . In Fig. 8, for the optimum paths  $\pi_s$  and  $\pi'_s$  we have:  
 $f_{\uparrow}(\pi_s \cdot \langle s, t \rangle) = 5 \neq f_{\uparrow}(\pi'_s \cdot \langle s, t \rangle) = 7$ .  
 Consequently,  $f_{\uparrow}$  doesn't satisfy this property ( $f_{\uparrow} \notin C4$ ).

### C. Peak amplitude of the relative intensity ( $f_{\max|\Delta I}$ )

Let  $f_{\max|\Delta I}$  be the function defined by:

$$f_{\max|\Delta I}(\pi_t = \langle t \rangle) = \begin{cases} 0, & \text{if } t \in S, \\ +\infty, & \text{otherwise.} \end{cases}$$

$$f_{\max|\Delta I}(\pi_t = \pi_s \cdot \langle s, t \rangle) = \max\{f_{\max|\Delta I}(\pi_s), |I(t) - I(R(\pi_s))|\}, \quad (8)$$

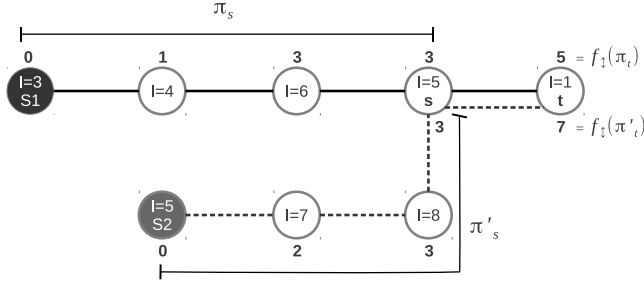


Fig. 8. Counterexample of the property C4 for the function  $f_{\downarrow}$ , where  $\pi_t = \pi_s \cdot \langle s, t \rangle$ ,  $\pi_s$ , and  $\pi'_s$  are optimum paths,  $\pi'_t = \pi'_s \cdot \langle s, t \rangle$  is a non-optimal path, and  $s_1, s_2$  are seeds.

where  $R(\pi_s)$  is the root pixel (origin) of the path  $\pi_s$ .

From the diagram, we have that  $f_{max|\Delta I|} \in C1 \setminus (C2 \cup C4)$ .

**Proof:**

- C1.  $f_{max|\Delta I|}(\pi_s) \leq f_{max|\Delta I|}(\pi_t)$ .  
 Suppose that this condition is true, then from (8) we have:  
 $f_{max|\Delta I|}(\pi_s) \leq \max\{f_{max|\Delta I|}(\pi_s), |I(t) - I(R(\pi_s))|\}$ ,  
 which is clearly right, because of the max function on the second term. Hence,  $f_{max|\Delta I|}(\pi_s)$  satisfies C1.
- C2.  $\pi_s$  is optimum.  
 In Fig. 9, for the optimum path  $\pi_t = \pi_s \cdot \langle s, t \rangle$ , we have that  $\pi_s$  is not optimum, since there is another path  $\pi'_s$  which offers a better cost than  $\pi_s$  to  $s$  (i.e.,  $f_{max|\Delta I|}(\pi'_s) = 7 \leq f_{max|\Delta I|}(\pi_s) = 8$ ). So,  $f_{max|\Delta I|}$  violates the condition C2 ( $f_{max|\Delta I|} \notin C2$ ).

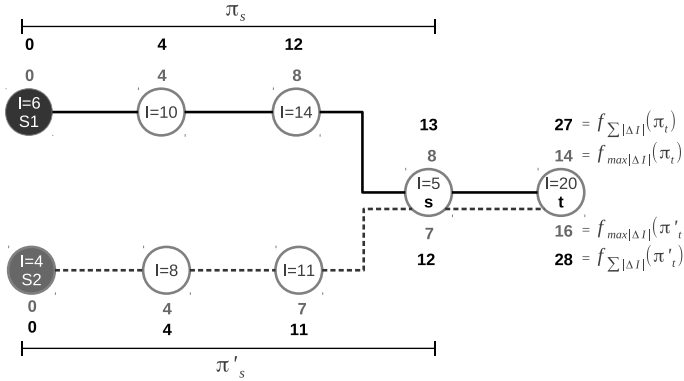


Fig. 9. Counterexamples of the second and third conditions (C2 and C3) for the functions  $f_{max|\Delta I|}$  and  $f_{\Sigma|\Delta I|}$ . The path  $\pi_t = \pi_s \cdot \langle s, t \rangle$  is an optimum path, and  $\pi'_t = \pi'_s \cdot \langle s, t \rangle$  is a non-optimal path, from seeds  $s_1$  and  $s_2$ , respectively.

- C3. For any optimum path  $\pi'_s$ ,  $f_{max|\Delta I|}(\pi'_s \cdot \langle s, t \rangle) = f_{max|\Delta I|}(\pi_t)$   
 We can see a counterexample in Fig. 9, where this condition is violated (i.e.,  $f_{max|\Delta I|}(\pi'_s \cdot \langle s, t \rangle) = 16 \neq f_{max|\Delta I|}(\pi_t) = 14$ ).
- C4. For any paths  $\pi_s$  and  $\pi'_s$  ending at  $s$ , if  $f_{max|\Delta I|}(\pi_s) = f_{max|\Delta I|}(\pi'_s)$ , then  $f_{max|\Delta I|}(\pi_s \cdot \langle s, t \rangle) = f_{max|\Delta I|}(\pi'_s \cdot \langle s, t \rangle)$ .

In Fig. 10, we can observe that:  $f_{max|\Delta I|}(\pi_s \cdot \langle s, t \rangle) = 14 \neq f_{max|\Delta I|}(\pi'_s \cdot \langle s, t \rangle) = 16$ . Hence,  $f_{max|\Delta I|}$  doesn't satisfy this property ( $f_{max|\Delta I|} \notin C4$ ).

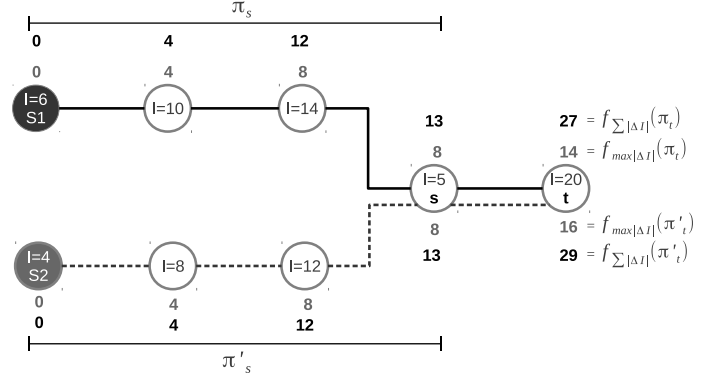


Fig. 10. Counterexamples of the condition C4 for the functions  $f_{max|\Delta I|}$  and  $f_{\Sigma|\Delta I|}$ . The paths  $\pi_t = \pi_s \cdot \langle s, t \rangle$ ,  $\pi_s$ , and  $\pi'_s$  are optimum,  $\pi'_t = \pi'_s \cdot \langle s, t \rangle$  is a non-optimal path, and  $s_1, s_2$  are seeds.

#### D. Sum of the absolute value of relative intensities

Let  $f_{\Sigma|\Delta I|}$  be the function defined by:

$$f_{\Sigma|\Delta I|}(\pi_t = \langle t \rangle) = \begin{cases} 0, & \text{if } t \in S, \\ +\infty, & \text{otherwise.} \end{cases}$$

$$f_{\Sigma|\Delta I|}(\pi_t = \pi_s \cdot \langle s, t \rangle) = f_{\Sigma|\Delta I|}(\pi_s) + |I(t) - I(R(\pi_s))|, \quad (9)$$

where  $R(\pi_s)$  is the root pixel (origin) of the path  $\pi_s$ .

**Proof:**

- C1.  $f_{\Sigma|\Delta I|}(\pi_s) \leq f_{\Sigma|\Delta I|}(\pi_t)$ .  
 Suppose that this condition is true, then from (9) we have:  
 $f_{\Sigma|\Delta I|}(\pi_s) \leq f_{\Sigma|\Delta I|}(\pi_s) + |I(t) - I(R(\pi_s))|$   
 As  $|I(t) - I(R(\pi_s))| \geq 0$  this relation is right, so the condition C1 is satisfied ( $f_{\Sigma|\Delta I|} \in C1$ ).
- C2.  $\pi_s$  is optimum.  
 In Fig. 9, for the optimum path  $\pi_t = \pi_s \cdot \langle s, t \rangle$ , we have that  $\pi_s$  is not optimum, since there is another path  $\pi'_s$  that has a better cost than  $\pi_s$  (i.e.,  $f_{\Sigma|\Delta I|}(\pi'_s) = 12 \leq f_{\Sigma|\Delta I|}(\pi_s) = 13$ ). So, this condition is not satisfied ( $f_{\Sigma|\Delta I|} \notin C2$ ).
- C3. For any optimum path  $\pi'_s$ ,  $f_{\Sigma|\Delta I|}(\pi'_s \cdot \langle s, t \rangle) = f_{\Sigma|\Delta I|}(\pi_t)$   
 The counterexample in Fig. 9 shows that this condition is violated (i.e.,  $f_{\Sigma|\Delta I|}(\pi'_s \cdot \langle s, t \rangle) = 28 \neq f_{\Sigma|\Delta I|}(\pi_t) = 27$ ).
- C4. For any paths  $\pi_s$  and  $\pi'_s$  ending at  $s$ , if  $f_{\Sigma|\Delta I|}(\pi_s) = f_{\Sigma|\Delta I|}(\pi'_s)$ , then  $f_{\Sigma|\Delta I|}(\pi_s \cdot \langle s, t \rangle) = f_{\Sigma|\Delta I|}(\pi'_s \cdot \langle s, t \rangle)$ .  
 In Fig. 10, for the optimum paths  $\pi_s$  and  $\pi'_s$  we have:  
 $f_{\Sigma|\Delta I|}(\pi_s \cdot \langle s, t \rangle) = 27 \neq f_{\Sigma|\Delta I|}(\pi'_s \cdot \langle s, t \rangle) = 29$ .  
 Hence,  $f_{\Sigma|\Delta I|}$  doesn't satisfy this property.

#### E. Function of intensity

Let  $f_I$  be the function defined by:

$$f_I(\pi_t = \langle t \rangle) = \begin{cases} 0, & \text{if } t \in S, \\ +\infty, & \text{otherwise.} \end{cases}$$

$$f_I(\pi_t = \pi_s \cdot \langle s, t \rangle) = I(t), \quad (10)$$

where  $S$  is a set of seeds.

**Proof:**

C1.  $f_I(\pi_s) \leq f_I(\pi_t)$ .

In Fig. 11 we have  $f_I(\pi_s) = 30 \not\leq f_I(\pi_t) = 20$ . So, this condition is violated.

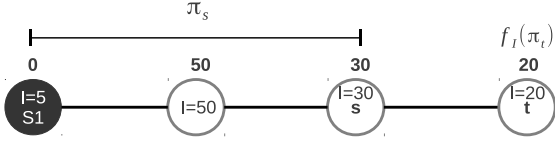


Fig. 11. Counterexample of the first condition (C1) for the function  $f_I$ , where  $\pi_t = \pi_s \cdot \langle s, t \rangle$  is an optimum path, and  $s_1$  is a seed.

C2.  $\pi_s$  is optimum.

This can be proved by contradiction as follows:

For  $\pi_t = \pi_s \cdot \langle s, t \rangle$  optimum, if  $\pi_s$  is not optimum then it is true that there exists an optimum path  $\pi'_s$ , such that  $f_I(\pi'_s) = I(s) < f_I(\pi_s) = I(s)$ , which is false, leading to a contradiction. Therefore, C2 is satisfied by  $f_I$ .

C3. For any optimum path  $\pi'_s$  ending at  $s$ ,  $f_I(\pi'_s \cdot \langle s, t \rangle) = f_I(\pi_t)$ .

This condition is also satisfied, since  $f_I(\pi'_s \cdot \langle s, t \rangle) = I(t)$  and  $f_I(\pi_t) = I(t)$ .

C4. For any paths  $\pi_s$  and  $\pi'_s$  ending at  $s$ , if  $f_I(\pi_s) = f_I(\pi'_s)$ , then  $f_I(\pi_s \cdot \langle s, t \rangle) = f_I(\pi'_s \cdot \langle s, t \rangle)$ .

This property is satisfied, because for any  $\pi_s$ ,  $f_I(\pi_s \cdot \langle s, t \rangle) = I(t)$ .

Thus, function  $f_I$  violates condition C1, and is not smooth.

#### F. Function lexicographic sum of $f_{\max}$

Let  $f_{\Sigma \max}^{lex}$  be the function defined by:

$$f_{\Sigma \max}^{lex}(\pi_t = \langle t \rangle) = \begin{cases} (0, 0), & \text{if } t \in S, \\ (+\infty, +\infty), & \text{otherwise.} \end{cases}$$

$$f_{\Sigma \max}^{lex}(\pi_t = \pi_s \cdot \langle s, t \rangle) = (f_{\Sigma \max}(\pi_t), f_{\max}(\pi_t)). \quad (11)$$

where  $S$  is a set of seeds.

This function  $f_{\Sigma \max}^{lex}$  gives pairs of values that should be compared according to the lexicographical order. The first component is the non-smooth function  $f_{\Sigma \max}$ , and the second is the smooth function  $f_{\max}$ .

**Proof:**

C1.  $f_{\Sigma \max}^{lex}(\pi_s) \leq f_{\Sigma \max}^{lex}(\pi_t)$ .

Suppose that this is true. Then, from (11) we have:

$$(f_{\Sigma \max}(\pi_s), f_{\max}(\pi_s)) \leq (f_{\Sigma \max}(\pi_t), f_{\max}(\pi_t)). \quad (12)$$

By comparing the ordered pairs we obtain:

$$f_{\Sigma \max}(\pi_s) < f_{\Sigma \max}(\pi_t), \quad (13)$$

or

$$f_{\Sigma \max}(\pi_s) = f_{\Sigma \max}(\pi_t) \text{ and } f_{\max}(\pi_s) \leq f_{\max}(\pi_t) \quad (14)$$

Since both  $f_{\Sigma \max}$  and  $f_{\max}$  satisfy condition C1 for  $w(s, t) \geq 0$  (see Section V-A), we may conclude that  $f_{\Sigma \max}^{lex} \in C1$  for  $w(s, t) \geq 0$ .

C2.  $\pi_s$  is optimum.

In Fig. 12, for the optimum path  $\pi_t = \pi_s \cdot \langle s, t \rangle$ , we have that  $\pi_s$  is not optimum, since there is another path  $\pi'_s$  that offers a better cost than  $\pi_s$  (i.e.,  $f_{\Sigma \max}^{lex}(\pi'_s) = (9, 7) \leq f_{\Sigma \max}^{lex}(\pi_s) = (10, 5)$ ). Therefore, this condition is not satisfied by  $f_{\Sigma \max}^{lex}$ .

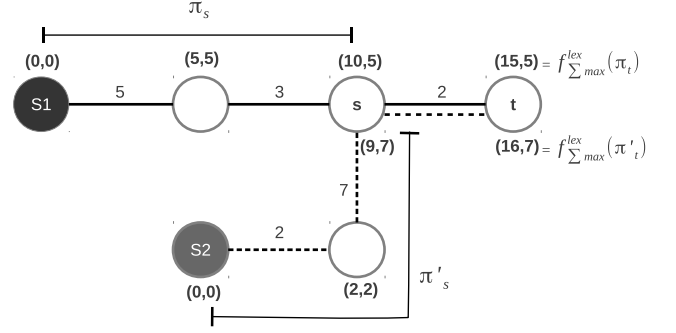


Fig. 12. Counterexample of the second and third conditions (C2 and C3) for the function  $f_{\Sigma \max}^{lex}$ , where  $\pi_t = \pi_s \cdot \langle s, t \rangle$  is an optimum path, and  $\pi'_t = \pi'_s \cdot \langle s, t \rangle$  is a non-optimal path, from the seeds  $s_1$  and  $s_2$ , respectively.

C3. For any optimum path  $\pi'_s$ ,  $f_{\Sigma \max}^{lex}(\pi'_s \cdot \langle s, t \rangle) = f_{\Sigma \max}^{lex}(\pi_t)$

In Fig. 12 we have

$$f_{\Sigma \max}^{lex}(\pi'_s \cdot \langle s, t \rangle) = (16, 7) \neq f_{\Sigma \max}^{lex}(\pi_t) = (15, 5).$$

Hence, this condition is violated by  $f_{\Sigma \max}^{lex}$ .

C4. For any paths  $\pi_s$  and  $\pi'_s$  ending at  $s$ , if  $f_{\Sigma \max}^{lex}(\pi_s) = f_{\Sigma \max}^{lex}(\pi'_s)$ , then  $f_{\Sigma \max}^{lex}(\pi_s \cdot \langle s, t \rangle) = f_{\Sigma \max}^{lex}(\pi'_s \cdot \langle s, t \rangle)$ . Given that  $f_{\Sigma \max}^{lex}(\pi_s) = f_{\Sigma \max}^{lex}(\pi'_s)$ , from (11) we obtain  $(f_{\Sigma \max}(\pi_s), f_{\max}(\pi_s)) = (f_{\Sigma \max}(\pi'_s), f_{\max}(\pi'_s))$ , which implies that:

$$f_{\Sigma \max}(\pi_s) = f_{\Sigma \max}(\pi'_s), \text{ and} \quad (15)$$

$$f_{\max}(\pi_s) = f_{\max}(\pi'_s). \quad (16)$$

Hence, by the definition of  $f_{\max}$  and from (16)

$$\begin{aligned} f_{\max}(\pi_s \cdot \langle s, t \rangle) &= \max\{f_{\max}(\pi'_s), \omega(s, t)\}, \\ f_{\max}(\pi_s \cdot \langle s, t \rangle) &= f_{\max}(\pi'_s \cdot \langle s, t \rangle) \end{aligned} \quad (17)$$

and from (4), (15), and (17) we have

$$\begin{aligned} f_{\Sigma \max}(\pi_s \cdot \langle s, t \rangle) &= f_{\Sigma \max}(\pi'_s) + f_{\max}(\pi'_s \cdot \langle s, t \rangle), \\ f_{\Sigma \max}(\pi_s \cdot \langle s, t \rangle) &= f_{\Sigma \max}(\pi'_s \cdot \langle s, t \rangle) \end{aligned} \quad (18)$$

So, from (17) and (18) we have

$$\begin{aligned} f_{\Sigma \max}^{lex}(\pi_s \cdot \langle s, t \rangle) &= (f_{\Sigma \max}(\pi'_s \cdot \langle s, t \rangle), f_{\max}(\pi'_s \cdot \langle s, t \rangle)), \\ f_{\Sigma \max}^{lex}(\pi_s \cdot \langle s, t \rangle) &= f_{\Sigma \max}^{lex}(\pi'_s \cdot \langle s, t \rangle) \end{aligned} \quad (19)$$

Therefore, the property C4 is satisfied by  $f_{\Sigma \max}^{lex}$ .

Thus, we have that the function  $f_{\Sigma \max}^{lex}$  is a non-smooth function that violates only the conditions C2 and C3. Note that  $f_{\Sigma \max}^{lex}$  does not violate C4, as opposed to  $f_{\Sigma \max}$ .